



Network Security Protocols

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March 2016



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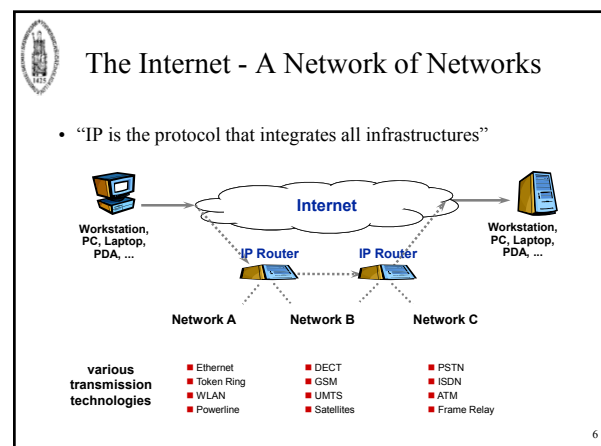
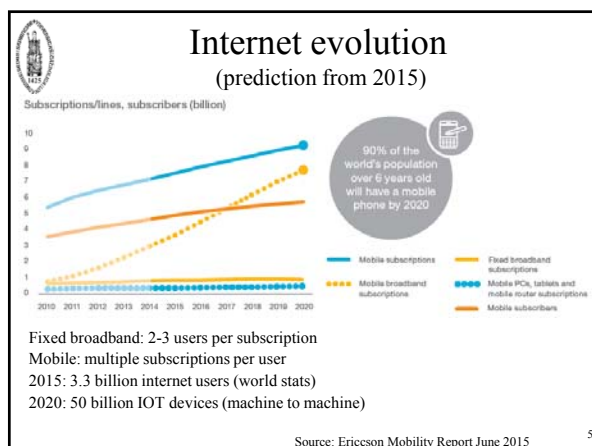
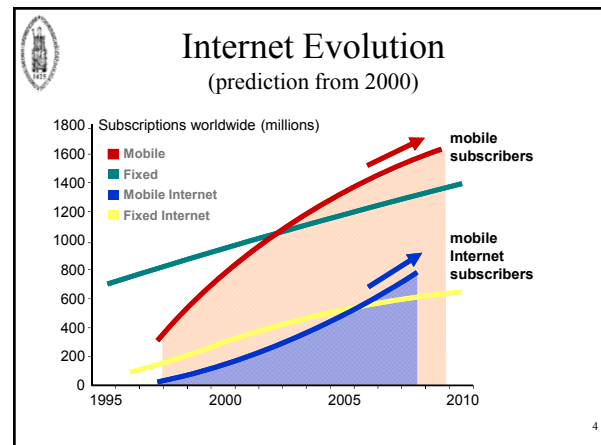
Goals

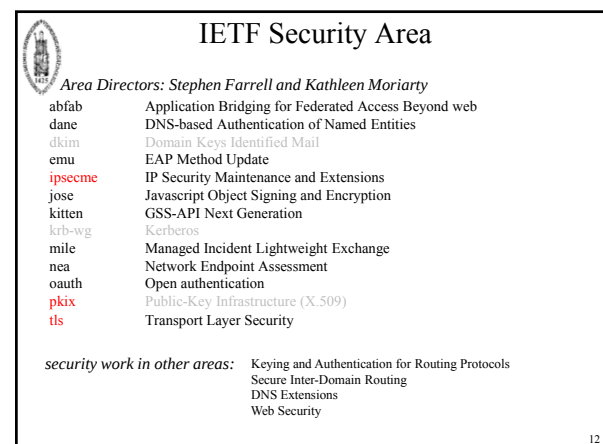
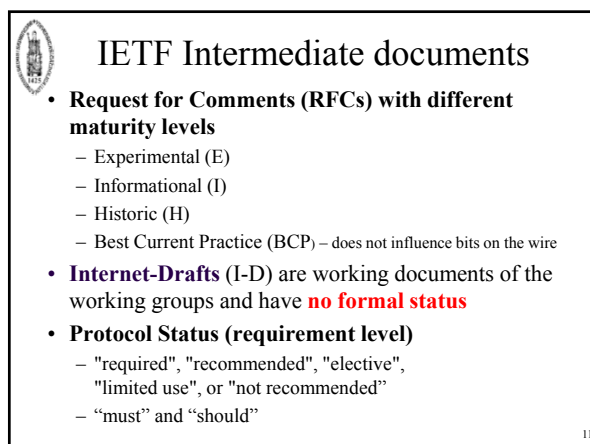
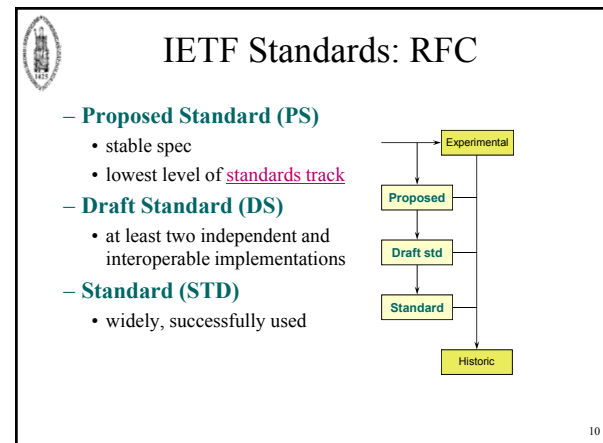
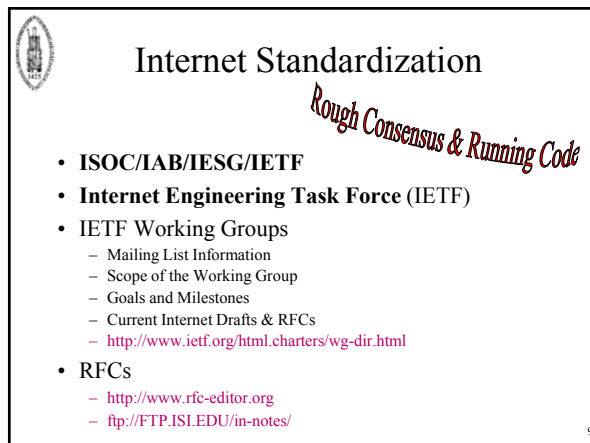
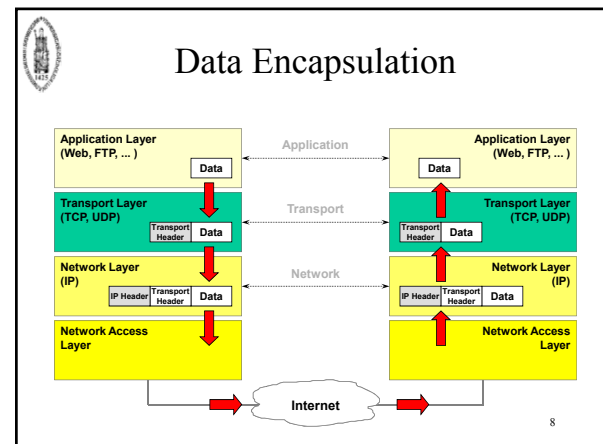
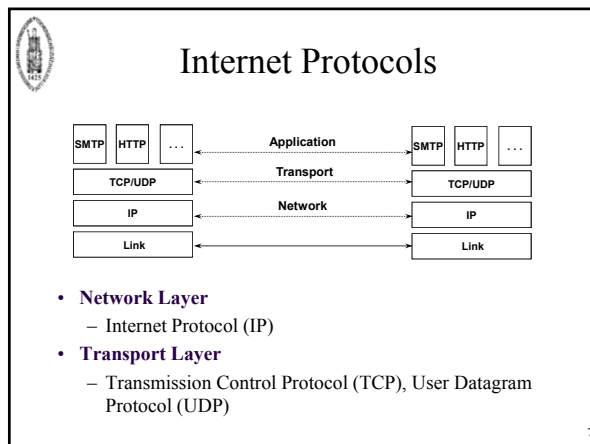
- Understanding how security can be added to the basic Internet protocols
- Understanding TLS and its limitations
- Understanding IPsec and its limitations



Outline

- Internet summary
- IETF process
- Basic principles
- Transport layer security
 - SSL / TLS
- Network layer security
 - IPSec, VPN, SSH



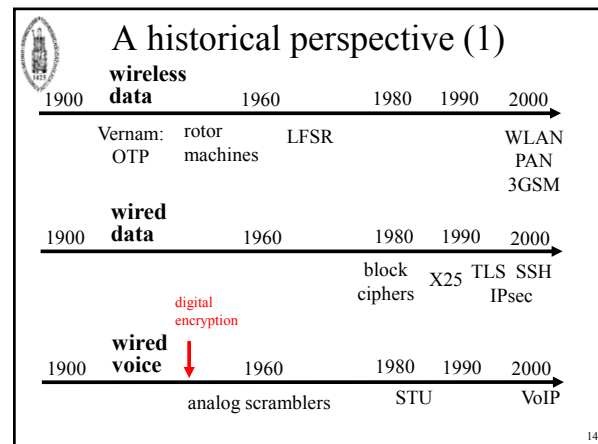


Communications insecurity

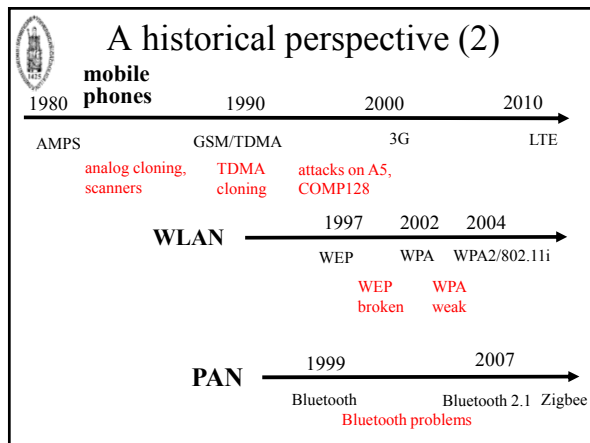
- architectural errors
 - wrong trust assumptions
 - default = no security
- protocol errors
 - unilateral entity authentication
 - weak entity authentication mechanism
 - downgrade attack
- modes of operation errors
 - no authenticated encryption
 - wrong use of crypto
- cryptographic errors
 - weak crypto
- implementation errors

range of wireless communication is often underestimated!

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Security Goals (started in ISO 7498-2)

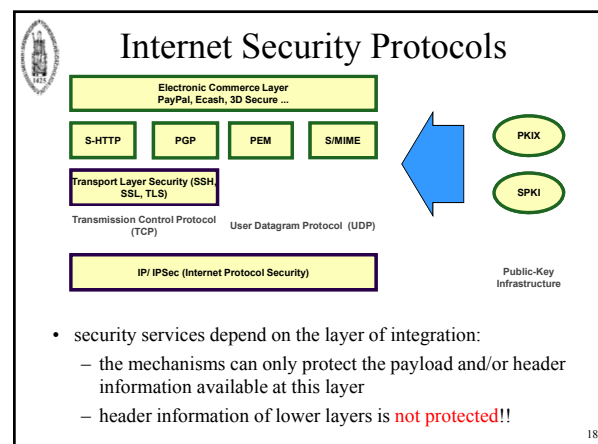
- confidentiality:
 - entities (anonymity)
 - data
 - traffic flow
- (unilateral or mutual) entity authentication
- data authentication (connection-less or connection-oriented): data origin authentication + data integrity
- access control
- non-repudiation of origin versus deniability

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Security Protocols & Services

- Cryptographic techniques:
 - symmetric encipherment
 - message authentication mechanisms
 - entity authentication mechanisms
 - key establishment mechanisms (e.g., combined with entity authentication)

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Security: at which layer?

- Application layer:
 - closer to user
 - more sophisticated/granular controls
 - end-to-end
 - but what about firewalls?
- Lower layer:
 - application independent
 - hide traffic data
 - but vulnerable in middle points
- Combine?

SP Architecture I: Encapsulation

- Bulk data: symmetric cryptography
- Authenticated encryption: best choice is to authenticate the ciphertext

SP Architecture II: Session (Association) Establishment

Algorithm Selection

"a la carte"

- each algorithm (encryption, integrity protection, pseudo-random function, Diffie-Hellman group, etc.) is negotiated independently
- less compact to encode
- more flexible
- e.g., IKEv1

"suite"

- all parameters are encoded into a single suite number; negotiation consists of offering one or more suites and having the other side choose
- simpler and more compact to encode
- potentially exponential number of suites
- less flexible
- e.g., TLS and IKEv2

Transport layer security

SSL / TLS

SSL/TLS Protocols

- connection-oriented data confidentiality and integrity, and optional client and server authentication.

Transport Layer Security Protocols

- IETF Working Group: **Transport Layer Security (tls)**
 - RFC 2246 (PS), 01/99
- transparent secure channels independent of the respective application.
- available protocols:
 - Secure Shell (SSH)*, SSH Ltd.
 - Secure Sockets Layer (SSL)*, Netscape
 - Transport Layer Security (TLS)*, IETF

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SSL / TLS

- Mainly in context of WWW security, i.e., to secure the HyperText Transfer Protocol (HTTP)
- TLS: security at the transport layer
 - can be used (and is intended) for other applications too (IMAP, telnet, ftp, ...)
 - end-to-end secure channel, but nothing more...
 - data is only protected during communication
 - no non-repudiation!

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Other WWW security protocols

- PCT: Microsoft's alternative to SSL
- S-HTTP: S/MIME-like protocol
- SET: e-payment protocol for credit card transactions
- XML-Signature: PKCS#7-based signature on XML documents
- ...

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SSL/TLS

- "Secure Sockets Layer" (Netscape)
 - SSL 2.0 (1995): security flaws!
 - SSL 3.0 (1996): still widely used - not interoperable with TLS 1.0
- "Transport Layer Security" (IETF)
 - TLS 1.0 (01/99) adopted SSL 3.0 with minor changes - RFC 2246 - default DSA/3DES
 - TLS 1.1 (4/2006) - RFC 4346 - default: RSA/3DES; several fixes for padding oracle and timing attacks (explicit IV for CBC)
 - TLS 1.2 (8/2008) - RFC 5246
 - replaces MD5 and SHA-1 by SHA-256 (SHA-1 still in a few places)
 - add AES ciphersuites (but still supports RC4!)
 - add support for authenticated encryption: GCM and CCM
 - RFC 5176 (2/2011) removes backward compatibility with SSL 2.0
 - Currently 314 ciphersuites!
 - TLS 1.3 expected for 2016**

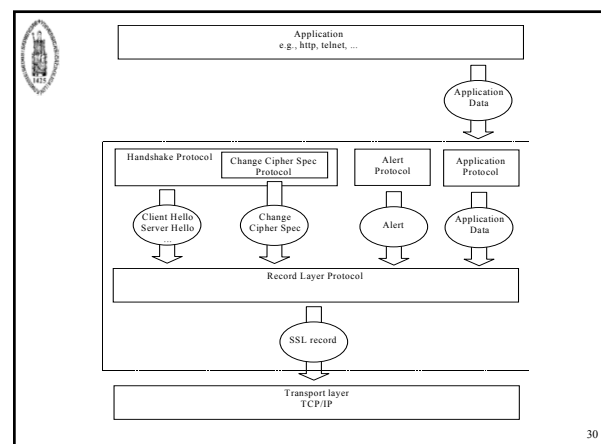
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SSL/TLS Deployment

more details: <https://www.trustworthyinternet.org/ssl-pulse/>

TLS 1.1 and 1.2 deployment very slow (about 25% of servers in Feb. 14); boost in Nov. 2013 (new attacks + Snowden revelations)

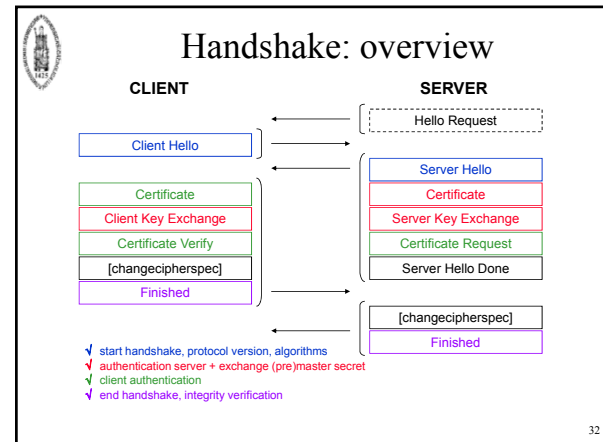
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SSL/TLS in more detail

- “Record layer” protocol
 - fragmentation
 - compression (not in practice) – will be removed in 1.3
 - cryptographic security:
 - encryption → data confidentiality
 - MAC → data authentication [no digital signatures!]
- “Handshake” protocol
 - negotiation of cryptographic algorithms
 - client and server authentication
 - establish cryptographic keys (master key and derived key for encryption and MAC algorithm)
 - key confirmation

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TLS 1.2 Data Encapsulation Options

Integrity			
key size	144	160	256
algorithm options	HMAC-MD5	HMAC-SHA	HMAC-SHA256

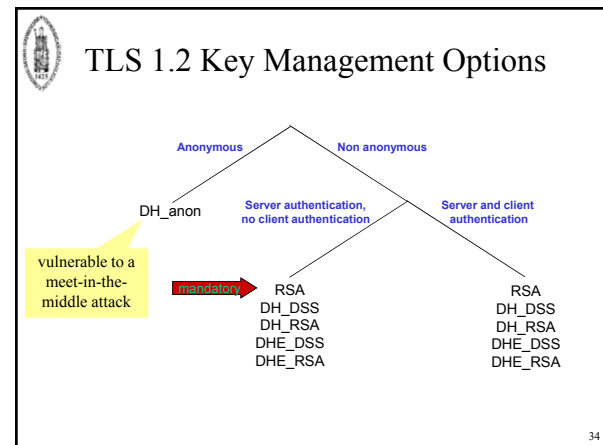
mandatory

Confidentiality					
key size	40	56	128	168	256
algorithm options	RC4_40 RC2_CBC_40 DES_CBC_40	DES_CBC	RC4_128 IDEA_CBC AES_CBC	3DES_EDE_CBC	AES_CBC

mandatory

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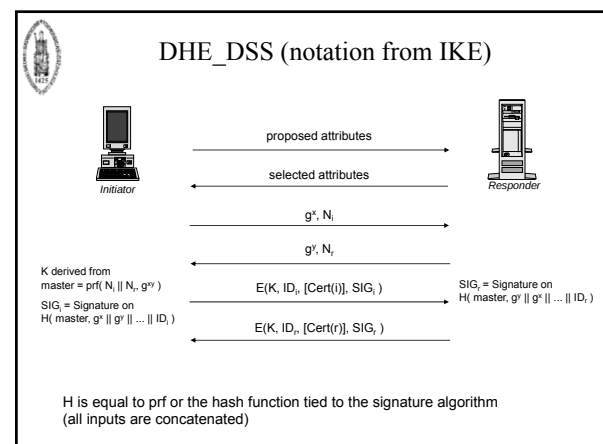
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Forward secrecy

- Default algorithm in TLS 1.2 is RSA (better performance, at least for RSA-1024)
 - no forward secrecy: compromise of private server key results in compromise of **all past** sessions
- DH-DSS and DH-DSA: same problem
- DHE-DSS and DHE-DSA: Ephemeral Diffie-Hellman keys leads to forward secrecy
 - For performance reasons: switch to a 256-bit Elliptic Curve (e.g. Google in November 2013)



SSL/TLS: security services

SSL/TLS only provides:

- entity authentication
- data confidentiality
- data authentication

SSL/TLS does not provide:

- non-repudiation
- unobservability (identity privacy)
- protection against traffic analysis
- secure many-to-many communications (multicast)
- security of the end-points (but relies on it!)

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SSL/TLS: security analysis

Detailed analysis and security reductions (“proofs”):

- Handshake protocol: most unaltered TLS ciphersuites form a secure channel (authenticated and confidential channel establishment)
- Record layer protocol: Authenticated Encryption well understood (but badly implemented)
- miTLS: validated reference implementation

Current analysis does not take into account the full complexity

- Cipher suites: negotiation, renegotiation, reuse of master key over multiple suites
- Cross protocol attacks
- Fragmentation
- Compression brings security problems
- Timing attacks

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TLS overview [Stebila'14]

Crypto primitives	Ciphersuite details	Protocol "Framework"	Libraries	Applications
RSA, DSA, ECDSA	Data structures	Alerts and errors	OpenSSL	Web browsers
DH, EC-DH	Key derivation	Certification/re-vocation	GnuTLS	Web servers
HMAC	Encryption modes and IVs	(Re-)Negotiation	SChannel	Application SDKs
MD5, SHA-1, SHA-2	Padding	Session Resumption	Java JSSE0	Certificates
DES, 3DES, RC4, AES	Compression	Key reuse		

Theoretical analysis

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TLS attack overview [Stebila'14]

Attacks shown:

- Rizzo & Duong "BEAST" attack
- Lucky 13
- Bleichenbacher RSA PKCSv1
- Debian OpenSSL entropy bug
- Rizzo & Duong "CRIME" attack
- RC4 biases
- Cross-protocol DV/COPI TLS attack
- Ray & Dispensa renegotiation attack
- SSL 2.0 downgrade
- Poor certificate validation
- CA breaches
- Goldberg & Wagner Netscape PRNG attack

TLS attacks (1)

- Renegotiation attack (2009)**
 - allows injection of data; patched by RFC 5746
- Version rollback attacks (2011)**
 - exploits false start feature (introduced to improve performance)
- CRIME and BREACH attacks (2013)**
 - recovery of cookies when *data compression* is used
 - all TLS versions are vulnerable
- Truncation attack (2013)**
 - suppress logout by injecting an unencrypted TCP FIN message
- Heartbleed (2014)**
 - Buffer over-read in OpenSSL implementation
- Poodle (2014)** Padding Oracle On Downgraded Legacy Encryption
 - Man-in-the middle that exploits downgrade to SSL 3.0
- Logjam (2015):** down negotiation to 512-DLOG that can be broken in real time
- SLOTH (2016):** TLS 1.2 allow use of pure MD5 – down negotiation
- DROWN (2016):** crossprotocol attack on SSLv2

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TLS attacks (2)

Padding oracle and timing attacks

- RSA**
 - [Bleichenbacher 98] PKCS #1v1.5 – 1 million chosen ciphertexts (in practice 200,000);
 - [Klimar+ 03] 40% improvement
 - [Bardou+ 12]: reduced to about 10,000 chosen ciphertexts
 - timing attack [Kocher'95], [Boneh-Brumley'03]
 - [Meyer+14] new Bleichenbacher attacks, even on TLS 1.3 draft
- CBC (IV and padding)**
 - padding [Rogaway], [Vaudenay 02], [Canvel+ 03]: password recovery
 - BEAST attack [Rizzo-Duon 11]: exploits IV issues - patched from TLS 1.1 onwards
 - Lucky 13 [AlFardan-Paterson'13]: timing attack on CBC padding
- Cryptographic attacks**
 - Weak random number generators: Netscape, Debian, embedded devices...
 - Exhaustive key search: 40-bit and 56-bit keys
 - Cross-protocol attack: elliptic curve parameters can be read as DH-prime
 - Biases in RC4 (re-introduced to 50% of web in Feb. 2013 to stop BEAST attack) [AlFardan+ 13] [Isobe+ 13]

More attacks and details: <https://mitsl.org/pages/attacks>

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TLS problems

- many PKI issues: revocation, root keys, fake certificates, certificate parsing,...
- web spoofing and phishing
- what if the user does not know that a particular website has to use SSL/TLS (solution HSTS – **HTTP Strict Transport Security (HSTS)**: mandate that you interact with particular servers using https/TLS only)
- traffic analysis:
 - length of ciphertext might reveal useful info
 - time to retrieve a page indicates whether it has been retrieved before

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TLS Renegotiation attack [Marsh Ray Nov.09]

- Cipher suite can be renegotiated dynamically throughout the session
 - negotiation and renegotiation look the same
- Person-In-The-Middle can inject (plaintext) traffic in a protected session as if it came from a client
- Fix: TLS renegotiation indication extension
RFC 5746 – Feb. '10
(84% deployment in Jan. '14)

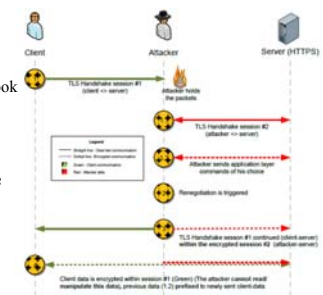


Figure: L. O'Connor

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Implementation attacks

Debian-OpenSSL incident [13 May 2008]
<https://cseweb.ucsd.edu/~hovav/dist/debiankey.pdf>

- Weak key generation:
 - only 32K keys
 - easy to generate all private keys
 - collisions
- Between 13-17 May 2008
 - 280 bad keys out of 40K (0.6%)
- Revocation problematic

IT'S NOT EASY TO GET THESE DONE

IN THE RUSH TO CLEAN UP THE DEBAY-CORRIGAN PIASCO, A NUMBER OF OTHER MARCO SECURITY HOLDS HAVE BEEN UNKNOWN:

AFFECTED SYSTEM	SECURITY PROBLEM
FEDORA CORE <i>(Fedora 10, 11, 12)</i>	VULNERABLE TO CERTAIN DECODER RINGS
MANDRIVO <i>(Mandriva 2009)</i>	GIVES ROOT ACCESS IF ASKED IN WERN VIO
GENTOO	VULNERABLE TO FLATTERY
OLPC OS	VULNERABLE TO TEST GOLDMINE'S POWERBOOK
SLACKWARE	GIVES ROOT ACCESS IF USER SAYS ENGLISH WORD FOR TREND
UBUNTU	TURNS OFF CUSTODY IS ACTUALLY NOT WANG WIVES WITH A FEW CUSTODY THIEVES

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TLS certificate "NULL" issue

- [Moxie Marlinspike 09] Black Hat
 - browsers may accept bogus SSL certs
 - CAs may sign malicious certs
- certificate for www.paypal.com m0.kuleuven.be will be issued if the request comes from a kuleuven.be admin
- response by PayPal: suspend Moxie's account
 - http://www.theregister.co.uk/2009/10/06/paypal_banishes_ssl_hacker/

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User authentication

First *authentication*, then *authorization* !

SSL/TLS client authentication:

- During handshake, client can digitally sign a specific message that depends on all relevant parameters of secure session with server
- Support by software devices, smart cards or USB tokens
- PKCS#12 key container provides software mobility
- rarely implemented

Usually another mechanism on top of SSL/TLS

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TLS 1.3

- Reduce the number of cipher suites:
 - only authenticated encryption with associated data (AEAD): AES-GCM, AES-CCM, ARIA-GCM, Camellia-GCM, ChaCha/Poly1305
 - only perfect forward secrecy (still RSA for signatures)
 - no custom DH groups
 - Forbid renegotiation but keep resumption with tickets
 - Improve privacy: encrypt more of the handshake
 - Improve latency: target: 1-RTT handshake for naive clients but **0-RTT handshake** for repeat connections
- Backward compatibility remains very important because of huge installed base

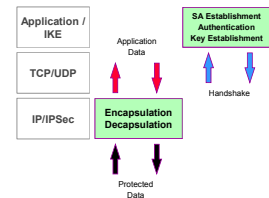
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Network layer security

IPsec, VPN, SSH

IP Security Protocols

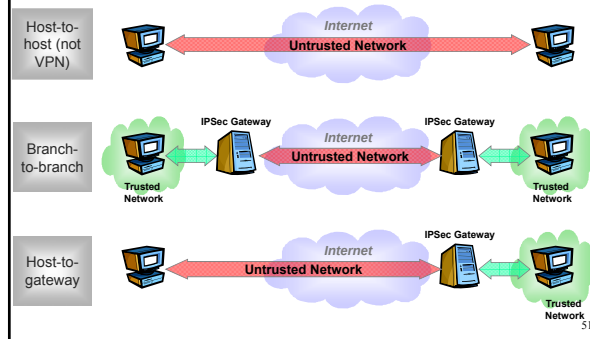
- IETF Working Group:
IP Security Protocol (ipsec)
Security Architecture for the Internet Protocol
 - RFC 2401 (PS), 11/98
- **IP Authentication Header (AH)**
 - RFC 2402 (PS), 11/98
- **IP Encapsulating Security Payload (ESP)**
 - RFC 2406 (PS), 11/98
- **Internet Key Exchange (IKE)**
 - RFC 2409 (PS), 11/98
 - Application layer protocol for negotiation of Security Associations (SA) and Key Establishment



- Large and complex..... (48 documents)
- Mandatory for IPv6, optional for IPv4

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IPsec VPN models: Hosts and Security Gateways



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IPsec - Security services

- Access control
- Connectionless integrity
- Data origin authentication
- Rejection of replayed packets (a form of partial sequence integrity)
- Confidentiality
- Limited traffic flow confidentiality

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IPsec - Concepts

- Security features are added as extension headers that follow the main IP header
 - Authentication header (AH)
 - Encapsulating Security Payload (ESP) header
- Security Association (SA)
 - Security Parameter Index (SPI)
 - IP destination address
 - Security Protocol Identifier (AH or ESP)

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IPsec - Parameters

- sequence number counter
- sequence counter overflow
- anti-replay window
- AH info (algorithm, keys, lifetimes, ...)
- ESP info (algorithms, keys, IVs, lifetimes, ...)
- lifetime
- IPsec protocol mode (tunnel or transport)
- path MTU (maximum transmission unit)

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IKE Algorithm Selection Mandatory Algorithms		
Algorithm Type	IKE v1	IKE v2
Payload Encryption	DES-CBC	AES-128-CBC
Payload Integrity	HMAC-MD5 HMAC-SHA1	HMAC-SHA1
DH Group	768 Bit	1536 Bit
Transfer Type 1 (Encryption)	ENCR_DES_CBC	ENCR_AES_128_CBC
Transfer Type 2 (PRF)	PRF_HMAC_SHA1 [RFC2104]	PRF_HMAC_SHA1 [RFC2104]
Transfer Type 3 (Integrity)	AUTH_HMAC_SHA1_96 [RFC2404]	AUTH_HMAC_SHA1_96 [RFC2404]

Source: draft-ietf-ipsec-ikev2-algorithms-00.txt, May 2003

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IPsec - Modes

- Transport (*host-to-host*)
 - ESP: encrypts and optionally authenticates IP payload, but not IP header
 - AH: authenticates IP payload and selected portions of IP header
- Tunnel (*between security gateways*)
 - after AH or ESP fields are added, the entire packet is treated as payload of new outer IP packet with new outer header
 - used for VPN

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IPsec - AH Transport mode

- Security Parameters Index: identifies SA
- Sequence number: anti-replay
- Integrity Check Value: data authentication using HMAC-SHA-1-96 or HMAC-MD5-96

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IPsec - AH Tunnel mode

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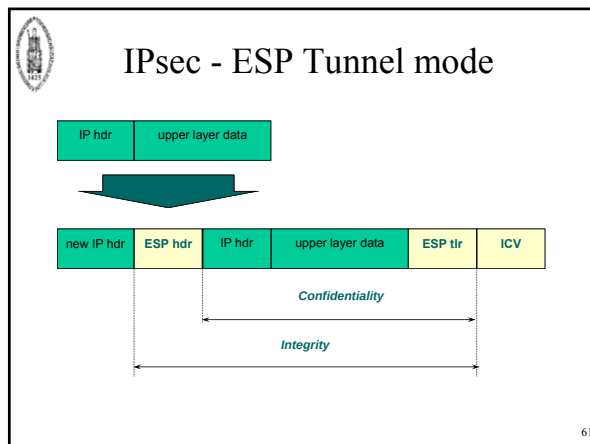
IPsec - ESP header

- Security Parameters Index: identifies SA
- Sequence number: anti-replay
- Encrypted payload data: data confidentiality using DES, 3DES, RC5, IDEA, CAST, Blowfish
- Padding: required by encryption algorithm (additional padding to provide traffic flow confidentiality)
- Integrity Check Value : data authentication using HMAC-SHA-1-96 or HMAC-MD5-96

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IPsec - ESP Transport mode

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IPsec: Key management

- RFCs 2407, 2408, and 2409
- Manual
- Automated
 - procedure / framework
 - Internet Security Association and Key Management Protocol (ISAKMP), RFC 2408 (PS)
 - key exchange mechanism: Internet Key Exchange (IKE)
 - Oakley: DH + cookie mechanism to thwart clogging attacks
 - SKEME

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IPsec: Key management

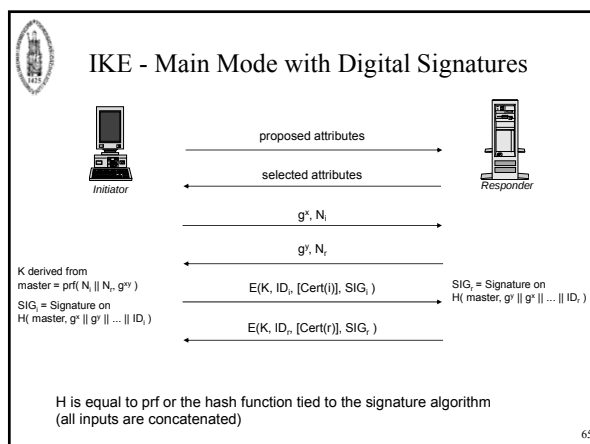
- IKE defines 5 exchanges
 - Phase 1: establish a secure channel
 - Main mode
 - Aggressive mode
 - Phase 2: negotiate IPSEC security association
 - Quick mode (only hashes, PRFs)
 - Informational exchanges: status, new DH group
- based on 5 generic exchanges defined in ISAKMP
- cookies for anti-clogging

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IPsec: Key management

- protection suite (negotiated)
 - encryption algorithm
 - hash algorithm
 - authentication method:
 - preshared keys, DSA, RSA, encrypted nonces
 - Diffie Hellman group: 5 possibilities

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IKE - Main Mode with Digital Signatures

- mutual entity authentication
- mutual implicit and explicit key authentication
- mutual key confirmation
- joint key control
- identity protection
- freshness of keying material
- perfect forward secrecy of keying material
- non-repudiation of communication
- cryptographic algorithm negotiation

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IKE v2 - RFC Dec 2005

- IKEv1 implementations incorporate additional functionality including features for NAT traversal, legacy authentication, and remote address acquisition, not documented in the base documents
- Goals of the IKEv2 specification include
 - to specify all that functionality in a single document
 - to simplify and improve the protocol, and to fix various problems in IKEv1 that had been found through deployment or analysis
- IKEv2 preserves most of the IKEv1 features while redesigning the protocol for efficiency, security, robustness, and flexibility

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IKE v2 Initial Handshake (1/2)

- Alice and Bob negotiate cryptographic algorithms, mutually authenticate, and establish a session key, creating an IKE-SA
- Usually consists of two request/response pairs
 - The first pair negotiates cryptographic algorithms and does a Diffie-Hellman exchange
 - The second pair is encrypted and integrity protected with keys based on the Diffie-Hellman exchange

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IKE v2 Initial Handshake (2/2)

- Second exchange
 - divulge identities
 - prove identities using an integrity check based on the secret associated with their identity (private key or shared secret key) and the contents of the first pair of messages in the exchange
 - establish a first IPsec SA (“child-SA”) is during the initial IKE-SA creation

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IPsec Overview

- much better than previous alternatives
- IPsec documents hard to read
- committee design: too complex
 - ESP in Tunnel mode with authenticated encryption probably sufficient
 - simplify key management
 - clarify cryptographic requirements
- ...and thus difficult to implement (securely)
- **avoid encryption without data authentication**

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VPN?

- Virtual Private Network
- Connects a private network over a public network.
- Connection is secured by tunneling protocols.
- The nature of the public network is irrelevant to the user.
- It appears as if the data is being sent over the private network
 - remote user access over the Internet
 - connecting networks over the Internet
 - connection computers over an intranet

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Concluding comments

- IPsec is really transparent, SSL/TLS only conceptually, but not really in practice
- SSH, PGP: stand-alone applications, immediately and easy to deploy and use
- Network security: solved in principle but
 - many implementation issues
 - complexity creates security weaknesses
- Application and end point security: more is needed!

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More information (1)

- William Stallings, *Cryptography and Network Security - Principles and Practice*, Fifth Edition, 2010
- N. Doraswamy, D. Harkins, *IPSec (2nd Edition)*, Prentice Hall, 2003 (outdated)
- Erik Rescorla, *SSL and TLS: Designing and Building Secure Systems*, Addison-Wesley, 2000.
- IETF web site: www.ietf.org
 - e.g., IETF-TLS Working Group
<http://www.ietf.org/html.charters/tls-charter.html>

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More information (2)

- Jon C. Snader, *VPNs Illustrated: Tunnels, VPNs, and IPsec*, Addison-Wesley, 2005
- Sheila Frankel, *Demystifying the IPsec Puzzle*, Artech House Computer Security Series, 2001
- Anup Gosh, *E-Commerce Security, Weak Links, Best Defenses*, Wiley, 1998
- Rolf Oppliger, *Security Technologies for the World Wide Web*, Artech House Computer Security Series 1999
- W3C Security (incl WWW Security FAQ)
<http://www.w3.org/Security/>

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